Digital Logic Structures

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Boolean Algebra

The Transistor

Logic Gates

Combinational Circuits

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Boolean Algebra

 A mathematical discipline used in the analysis and design of digital circuits

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- It uses logical variables and operations
- Variables:
 - 1 (TRUE)
 - 0 (FALSE)
- Operations:
 - A AND $B = A \cdot B$
 - A OR B = A + B
 - NOT A = \overline{A}

Identities of Boolean Algebra

Basic Postulates			
$\mathbf{A} \cdot \mathbf{B} = \mathbf{B} \cdot \mathbf{A}$	A + B = B + A	Commutative Laws	
$\mathbf{A} \cdot (\mathbf{B} + \mathbf{C}) = (\mathbf{A} \cdot \mathbf{B}) + (\mathbf{A} \cdot \mathbf{C})$	$\mathbf{A} + (\mathbf{B} \cdot \mathbf{C}) = (\mathbf{A} + \mathbf{B}) \cdot (\mathbf{A} + \mathbf{C})$	Distributive Laws	
$1 \cdot \mathbf{A} = \mathbf{A}$	0 + A = A	Identity Elements	
$\mathbf{A} \cdot \overline{\mathbf{A}} = 0$	$A + \overline{A} = 1$	Inverse Elements	
Other Identities			
$0 \cdot \mathbf{A} = 0$	1 + A = 1		
$A \cdot A = A$	A + A = A		
$\mathbf{A} \cdot (\mathbf{B} \cdot \mathbf{C}) = (\mathbf{A} \cdot \mathbf{B}) \cdot \mathbf{C}$	A + (B + C) = (A + B) + C	Associative Laws	
$\overline{\mathbf{A} \cdot \mathbf{B}} = \overline{\mathbf{A}} + \overline{\mathbf{B}}$	$\overline{\mathbf{A} + \mathbf{B}} = \overline{\mathbf{A}} \cdot \overline{\mathbf{B}}$	DeMorgan's Theorem	

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The Transistor

A transistor is a semiconductor device used to amplify and switch electronic signals

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Logic Gates

Electronic circuits that produce output signals that are Boolean operations on their input signals

When the gate input changes, the output is instant unless a gate delay occurs as a result of signal propagation time

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Digital logic uses the following gates:

- AND
- OR
- NOT
- NAND
- NOR
- XOR

The NOT Gate (inverter)







 $\begin{array}{c|ccc}
A & B & A+B \\
\hline
0 & 0 & 0 \\
0 & 1 & 1 \\
1 & 0 & 1 \\
1 & 1 & 1 \\
\end{array}$



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The XOR Gate

Α	В	$\pmb{A} \oplus \pmb{B}$
0	0	0
0	1	1
1	0	1
1	1	0



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The NOR Gate

Α	В	$\overline{A+B}$
0	0	1
0	1	0
1	0	0
1	1	0



The AND Gate

Α	В	AB
0	0	0
0	1	0
1	0	0
1	1	1

How would you construct an AND gate?

The NAND gate

Α	В	AB
0	0	1
0	1	1
1	0	1
1	1	0



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Logic Gate Symbols

Assert a signal is to cause a signal line transition from its logically false (0) state to its logically true (1) state



DeMorgan's Law

$$AB = \overline{\overline{A} + \overline{B}}$$

$$A + B = \overline{\overline{A} \cdot \overline{B}}$$

Show the relationship between AND and OR gates. AND, OR, NAND, and NOR gates may have more than two inputs.

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Complete Sets of gates

Any Boolean function can be implemented using a subset of gates

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We need to identify a functionally complete set of gates

The following sets of gates are functionally complete

- AND, OR, NOT
- AND, NOT
- OR, NOT
- NAND
- NOR

Using NAND Gates



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Using NOR Gates



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Combinational Circuits

- An interconnected set of gates
- The output at any time is a function only of the input at that time

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- The output is determined by the input
- Examples
 - Adder
 - Decoder
 - MUX

Design of Combinational Circuits

- Problem statement
- Formal specification
 - Truth table
 - Boolean equation
 - Graphical symbols: hardware description language (eg VHDL)

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- Minimize
- Implementation

Design of Combinational Circuits

Example

A car key alert should be turned on under the following circumstances:

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- The key is in and the door is open;
- The key is not in and the light is on.

The system has three inputs:

- L the lights are on
- ► K the key is in
- D the door is open

The system has one output:

B - the buzzer will be activated

Description to Truth Table

From the description we can obtain the truth table

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L	Κ	D	В
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	1
1	0	0	1
1	0	1	1
1	1	0	0
1	1	1	1

Sum of Products

The Sum of Products (SOP) has one term for each 1 in the output

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 $B = \overline{L}KD + L\overline{K}\overline{D} + L\overline{K}D + LKD$

Product of Sums

The Product of Sums (POS) has one term for each 0 in the output



 $(\overline{L} + \overline{K} + \overline{D}) \cdot (\overline{L} + \overline{K} + D) \cdot (\overline{L} + K + \overline{D}) \cdot (L + K + \overline{D})$

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Definitions

Definition (Literal)

A variable X has two literals X and \overline{X} .

Definition (Product term)

A logical product where each variable is represented by at most one literal is a **product** or a **product term** or a **term**. A term can be a single literal.

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Definition (Sum-of-products)

A logical sum of product terms forms a sum-of-products expression (SOP).

Definitions

Definition (Minterm)

A **minterm** is a logical product of *n* literals where each variable occurs as exactly one literal

Definition (Canonical SOP)

A **canonical SOP** is a logical sum of minterms, where all minterms are different. Also called canonical disjunctive form or minterm expansion.

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Problem

- List the minterms; $B(L, K, D) = \Sigma m(3, 4, 5, 7)$
- ► Write the canonical SOP for the car buzzer problem; $B(L, K, D) = \overline{L}KD + L\overline{KD} + L\overline{KD} + L\overline{KD}$

- draw the circuit;
- how can it be minimized?
- show that $L\overline{K} + KD + LD = L\overline{K} + KD$

Simplification of a Boolean Expression

A Boolean function can be realized in either SOP or POS form

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However we need a simpler Boolean expression that uses fewer gates:

- Algebraic simplification
- Karnaugh maps
- Quine-McCluskey tables

Algebraic Simplification

- Involves the use of boolean algebra identities
- To reduce the Boolean expression to one with fewer elements
- Or even simple observation, for example:

$$F = \overline{A}B + B\overline{C}$$

Can be simplified as:

$$F = B(\overline{A} + \overline{C})$$

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