

Dynamic Admission Control For A Bandwidth Broker Chenguang Gao, John DeDourek, Przemyslaw Pochec Faculty of Computer Science, University of New Brunswick Fredericton, NB, Canada



INTRODUCTION

• Multimedia and real-time applications require high quality services.

• Quality of Service (QoS) provides better service such as reduction of the number of dropped packets, delay, jitter, and out-of-order delivery.

Methodology

• Flow Generator - randomly generate EF flows and store them into a file

• NS2 DiffServ Script – provide the DiffServ, monitor packet dropping, read incoming flow file and bandwidth file

• Smart Admission Control - read flow files, inspect network load, predict and generate future threshold for EF traffic

Admission Control Algorithms •Static Admission Control: AT(0) = AT(1) = AT(2) = = AT(N) = Initial AT

•Dynamic Admission Control: $AT(N) = AT(N - 1)_{(optimal)}$ $\eta = \alpha \times UsedCapacity - \beta \times UnusedCapacity$ $-\gamma \times \text{RejectFlows}$ ($\alpha = \beta = \gamma = 1$)

	Flow	Data Source			
	Rate	start	stop	source	destination
flow 0	1000000	10	40	0	0
flow 1	2000000	15	35	1	1
flow 2	2000000	20	30	2	2
flow 3	2000000	25	35	3	3
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• The IETF proposed the Differentiated Services (DiffServ), which classifies flows with DiffServ code point (DSCP) and the Per-Hop behavior (PHB).

• A Bandwidth Broker (BB) manages the resources based on the Service Level Agreement (SLA) by controlling the network load and by accepting or rejecting bandwidth requests.

OBJECTIVE

•Design a scheme to provide dynamic bandwidth management in a DiffServ domain with a bandwidth broker.

•Simulate the proposed scheme in NS-2 and analyze results with respect to performance of the admitted streams, and with respect to the cost of unused reserved resources





Experiments and Results

flowLostEF.tr flowLostEE.tr



Overall Flowchart



Flow Generating Algorithm

algorithm --> generating EF flows for each pair of EF_node(\$c) initialize totol running time (sum(\$c)) initialize start_time for each flow (start(\$c,0)) initialize stop_time for each flow (stop(\$c,0))

set up RNG and its seed (it could be 0 or a constant

- set up and generate cbr_rate(\$c,0) with Uniform distribution
- set up and generate idle(\$c,0) with Uniform distribution
- set up start(\$c,0) with idle(\$c,0)
- set up and generate video_length (\$c,0) with Uniform distribution

add up \$idle(\$c,0)+ video_length(\$c,0) to sum(\$c) define stop(\$c,0) with \$sum(\$c)

initialize counter d to 1

while $\{$ sum(sc) < 660seconds $\}$

generate rate(\$c,\$d) with Uniform distribution generate idle(\$c,\$d) with Uniform distribution set up start (c, d) with [sidle(c, d) + stop(c, d-1)

generate video_length(\$c,\$d) with Uniform distribution add up to $sum(sc)+ sidle(sc,sd)+ svideo_length(sc,sd)$ to sum(sc)set up stop(\$c,\$d) with the new \$sum(\$c)

increase counter d



5.00 + 0

Conclusion

•The simulation shows improved QoS for the EF traffic with dynamic admission control (very few EF packets dropped). •Performance measured with the metric η is higher for the dynamic algorithm than for the static algorithm. •The proposed scheme successfully provides dynamic bandwidth management with a bandwidth broker.